

ДАА

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Hello, world!

...ОТ МИНАЛИЯ ПЪТ

$$\begin{aligned}
 \bullet T(n) &= n * T(n-1) + 1 = n(n-1)T(n-2) + 1 + 1 = \\
 &= n(n-1)T(n-2) + n + 1 = \\
 &= n(n-1)(n-2)T(n-3) + n(n-1) + n + 1 = \\
 &= n(n-1)(n-2)(n-3)T(n-4) + n(n-1)(n-2) + n(n-1) + n + 1 \\
 &= \dots = n(n-1)\dots 4 \cdot 3 \cdot 2 \cdot T(1) + n(n-1)\dots 4 \cdot 3 + n(n-1)\dots 4 + \dots + n(n-1) + n + 1 \\
 \bullet T(n) &= T(n-2) + 2 * \lg n
 \end{aligned}$$

$$\lim_{n \rightarrow \infty} \frac{\frac{1}{(n+1)!}}{\frac{1}{n!}} = \lim_{n \rightarrow \infty} \frac{n!}{(n+1)!} = \lim_{n \rightarrow \infty} \frac{1}{n+1} = 0 < 1 \Rightarrow \sum_{i=0}^{\infty} \frac{1}{i!} < \text{const}$$

$$\begin{aligned}
 &= n!T(1) + \frac{n!}{2!} + \frac{n!}{3!} + \frac{n!}{4!} + \dots + \frac{n!}{(n-2)!} + \frac{n!}{(n-1)!} + \frac{n!}{n!} = \\
 &= n! \Theta(1) + n! \sum_{i=2}^n \frac{1}{i!} < n! \Theta(1) + n! \cdot \text{const} \\
 &\geq n! \Theta(1) \Rightarrow T(n) = \Omega(n!) \Rightarrow T(n) = \Theta(n!)
 \end{aligned}$$

$$T(n) = T(n-2) + 2 \lg n =$$

$$= T(n-4) + 2 \lg(n-2) + 2 \lg n =$$

$$= T(n-6) + 2 \lg(n-4) + 2 \lg(n-2) + 2 \lg n$$

$$I_{\text{ка.}} = \dots = T(1) + 2 \lg 3 + \dots + 2 \lg(n-2) + 2 \lg n$$

n -независимо

$$= \Theta(1) + \underbrace{\lg 3 + \lg 3 + \lg 5 + \lg 5 + \dots + \lg(n-2) + \lg(n-2) + \lg n + \lg n}_{\lg(n-3) + \lg(n-2) + \lg(n-1) + \lg n} = S$$

$$T(n) = 2T\left(\frac{n}{2}\right) + \frac{n}{\lg n}$$

разобравшие?

$$S > \lg 2 + \lg 3 + \lg 4 + \lg 5 + \dots + \lg(n-3) + \lg(n-2) + \lg(n-1) + \lg n$$

$$\Leftrightarrow S > \lg(n!) \approx n \lg n \Rightarrow S = \Omega(n \lg n)$$

$$S < \lg 3 + \lg 4 + \lg 5 + 6 + \dots + \lg n + \lg(n+1)$$

$$S < \frac{\lg(n+1)!}{\lg 2} = \Theta((n+1) \lg(n+1)) = \Theta(n \lg n)$$

$$\Rightarrow S = O(n \lg n) \Rightarrow S = \Theta(n \lg n)$$

1. $T(n) = T(n - 1) + 1$

2. $T(n) = T(n - 2) + (2n)^{\wedge}$

3. $T(n) = 2T(n - 1) + 2^n$

4. $T(n) = T(n - 1) + T(n - 2)$

5. $T(n) = 2T(n - 1) - T(n - 2)$

6. $T(n) = 4T(n - 2) + 3^n$

7. $T(n) = 3T(n - 1) - 2T(n - 2) + n2^n$

8. $T(n) = \sum_{i=1}^{n-1} T(i) + 2^n$

9. $T(n) = \sum_{i=1}^{n-2} T(i) + 3^{\frac{n}{2}}$

10. $T(n) = 2T\left(\frac{n}{2}\right) + 1$

Handwritten notes for problem 6:
 $x^n = 2 \times x^{n-2} \Leftrightarrow x^{-2} = 0 \rightarrow \{2, 2\} \rightarrow \{-2, 2, 3\}$
 $T(n) = \frac{A \cdot (-2)^n + B \cdot 2^n + C \cdot 3^n}{O(2^n)} \approx O(3^n)$

Handwritten notes for problem 7:
 $x^2 - 3x + 2 = 0 \rightarrow \{2, 2\} \Rightarrow \{1, 2\} \Rightarrow \{1, 2, 2, 2\}$
 $T(n) = \sum_{i=1}^{n-1} T(i) + 2^n \Rightarrow T(n) = A \cdot 1^n + B \cdot 2^n + C \cdot n2^n + 2 \cdot n \cdot 2^n = O(n^2 \cdot 2^n)$

Handwritten notes for problem 10:
 $T(n) = \sum_{i=1}^{n-2} T(i) + 2^{\frac{n}{2}}$
 $T(n-1) = \sum_{i=1}^{n-2} T(i) + 2^{\frac{n-1}{2}}$
 $T(n) - T(n-1) = T(n-1) + 2^{\frac{n}{2}} - 2^{\frac{n-1}{2}} \Leftrightarrow T(n) = 2T(n-1) + \frac{1}{2} \cdot 2^{\frac{n}{2}}$
 $\Rightarrow T(n) = A \cdot 2^n + B \cdot n2^{\frac{n}{2}} = O(n \cdot 2^{\frac{n}{2}})$

$$T(n) = 2T\left(\frac{n}{2}\right) + 1$$

$$n = 2^m \quad T(n) = T(2^m) = S(m)$$

$$\lg n = m$$

$$\frac{n}{2} = \frac{2^m}{2} = 2^{m-1}$$

$$T\left(\frac{n}{2}\right) = T(2^{m-1}) = S(m-1)$$

$$\Rightarrow S(m) = 2S(m-1) + \underbrace{1}_{\{1\}} \cdot \underbrace{1^m}_{\{2\}}$$

$x=2 \Rightarrow \{2\}$

$$\Rightarrow \{1; 2\} \Rightarrow S(m) = A \cdot 1^m + B \cdot 2^m$$

$$T(n) = S(m) = \Theta(2^m) = \Theta(n)$$

$$T(n) = a \cdot T\left(\frac{n}{b}\right) + f(n)$$

$$k = \log_b a \quad (\text{патн. } f(n) \text{ с } n^k)$$

I ca. За $\varepsilon > 0$ $n^{k-\varepsilon} \nlessdot f(n) \Rightarrow T(n) \sim n^k$

II ca. $f(n) \sim n^k \cdot \lg^t n$ за $t \geq 0 \Rightarrow T(n) \sim n^k \cdot \lg^{t+1} n$

III ca. За $\varepsilon > 0$ $n^{k+\varepsilon} \nlessdot f(n) \Rightarrow T(n) \sim f(n)$

и за $c \in (0; 1)$ е вярно, че

за достатъчно големи n

$$a \cdot f\left(\frac{n}{b}\right) < c \cdot f(n)$$

$$1. T(n) = T\left(\frac{9n}{10}\right) + n$$

$$2. T(n) = 16T\left(\frac{n}{4}\right) + n^2$$

$$3. T(n) = 2T\left(\frac{n}{4}\right) + \sqrt{n}$$

$$4. T(n) = 4T\left(\frac{n}{2}\right) + n$$

$$5. T(n) = 3T\left(\frac{n}{2}\right) + n \rightarrow \text{Karatsuba}$$

$$6. T(n) = 3T\left(\frac{n}{4}\right) + n \lg n$$

$$7. T(n) = 8T\left(\frac{n}{2}\right) + n^3$$

$$8. T(n) = 2T\left(\frac{n}{2}\right) + n \rightarrow \text{MergeSort}$$

$$9. T(n) = 2T\left(\frac{n}{2}\right) + \lg n$$

$$10. T(n) = 2T\left(\frac{n}{2}\right) + n \lg n$$

$$11. T(n) = 5T\left(\frac{n}{2}\right) + n^2 \lg n$$

$$12. T(n) = 2T\left(\frac{n}{2}\right) + 1$$

$$13. T(n) = 2T(\sqrt{n}) + \lg n$$

$$4. k = \lg_2 4 = 2$$

Сравн. $n^k = n^2$ с $f(n) = n$

$$\stackrel{?}{=} \varepsilon > 0; n^{k-\varepsilon} \neq f(n)$$

$$- \exists \alpha \forall \varepsilon \in (0; 1) \text{ e } \text{всегда } 2 - \varepsilon \geq 1$$

$$\Leftrightarrow n^{k-\varepsilon} \neq n = f(n)$$

\Rightarrow no la. na MT

$$T(n) \approx n^2$$

$$3. T(n) = 2T\left(\frac{n}{4}\right) + \sqrt{n}$$

$$k = \log_4 2 = \frac{1}{2} \Rightarrow \text{сроби. } n^k = n^{\frac{1}{2}} \text{ c } f(n) = n^{\frac{1}{2}}$$

$$\Rightarrow f(n) = n^k \Rightarrow f(n) \asymp n^k \cdot \lg^n \text{ за } t=0$$

$$\Rightarrow \text{но ил. на МТ } T(n) \asymp \sqrt{n} \cdot \lg n$$

$$\log_a b = \frac{1}{\log_c a}$$

$$\text{Сроби. } n^k \text{ c } f(n) = 2^n$$

$$\text{За } \forall \epsilon > 0 \quad n^{k+\epsilon} \not\asymp 2^n$$

$$6. T(n) = 3T\left(\frac{n}{4}\right) + n \lg n$$

$$k = \log_4 3 \in (0; 1) \Rightarrow \text{сроби. } n^k \text{ c } f(n) = n \cdot \lg n$$

$$\rightarrow \text{за } \forall \epsilon \in (0, 1 - \log_4 3) \text{ e } \text{сроби. } n^{k+\epsilon} \not\asymp n \cdot \lg n$$

$$\text{Трчим } c \in (0; 1): \exists n_0 > 0: \forall n \geq n_0 \quad 3 \cdot \frac{n}{4} \cdot \lg \frac{n}{4} < c \cdot n \lg n$$

$$\Leftrightarrow \exists c: \frac{3}{4} \cdot n \cdot (\lg n - 2) < c \cdot n \cdot \lg n \quad | : n$$

$$\Leftrightarrow \exists c: \frac{3}{4} \cdot \lg n - \frac{3}{2} < c \cdot \lg n$$

$$\rightarrow \text{га, за } c = \frac{3}{4} \text{ за } \forall n > 0 \text{ e } \text{сроби}$$

но ил. на МТ

$$\Rightarrow T(n) \asymp n \lg n$$